# BALM: QoS-Aware Memory Bandwidth Partitioning for Multi-Socket Cloud Nodes

<u>David Gureya</u><sup>1,2</sup>, Vladimir Vlassov<sup>2</sup> and João Barreto<sup>1</sup>

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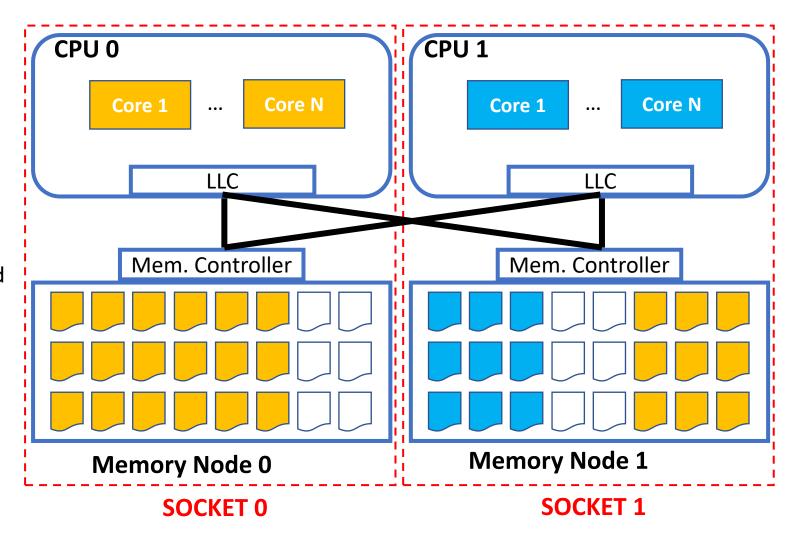
## Workload Consolidation (1/2)

- Widely used to improve resource utilization
  - Multiple workloads are consolidated on the same physical servers

## Workload Consolidation (2/2)

## Best-Effort Application (BEA)

- Throughput-oriented
   Bandwidth-intensive
   BEAs
- Place pages acrossMemory nodes

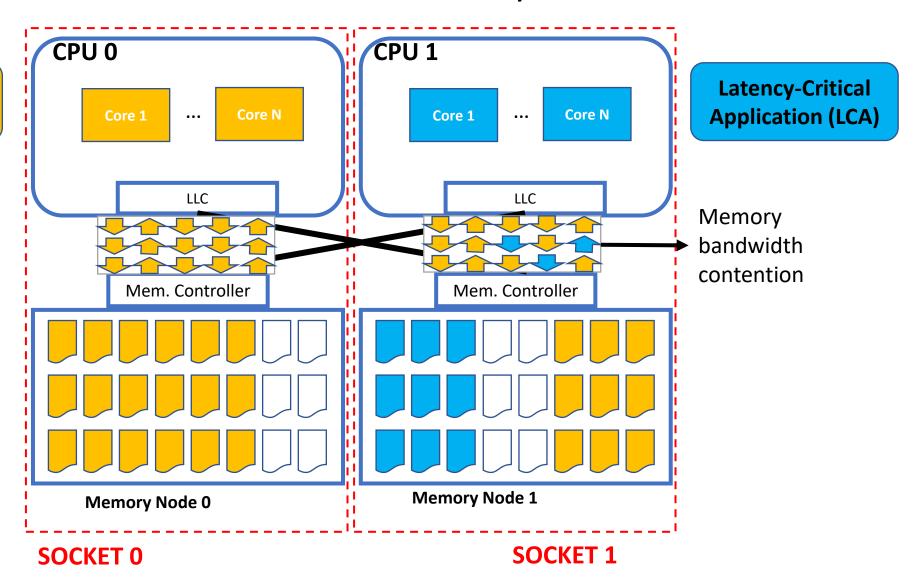


## Latency-Critical Application (LCA)

QoS requirements

### Challenge: Contention for Memory Bandwidth

Best-Effort
Application (BEA)



#### Problem: QoS-Aware Resource Allocation

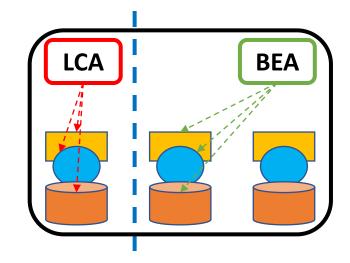
 Safeguard the SLO of LCAs while maximizing the throughput of the BEAs

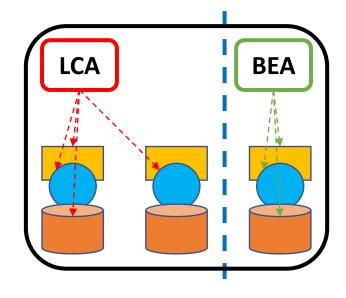
• **Dynamic problem**: resource usage by each application can change at any time

## Existing Solutions (1/2)

Solve the problem by partitioning resources

- Recent systems dynamically adjust partitions
  - Heracles [ISCA'15]
  - Parties [ASPLOS'19]
  - Clite [HPCA'20]
  - Caladan [OSDI'20]
- However, existing solutions are tailored to single-socket architectures only



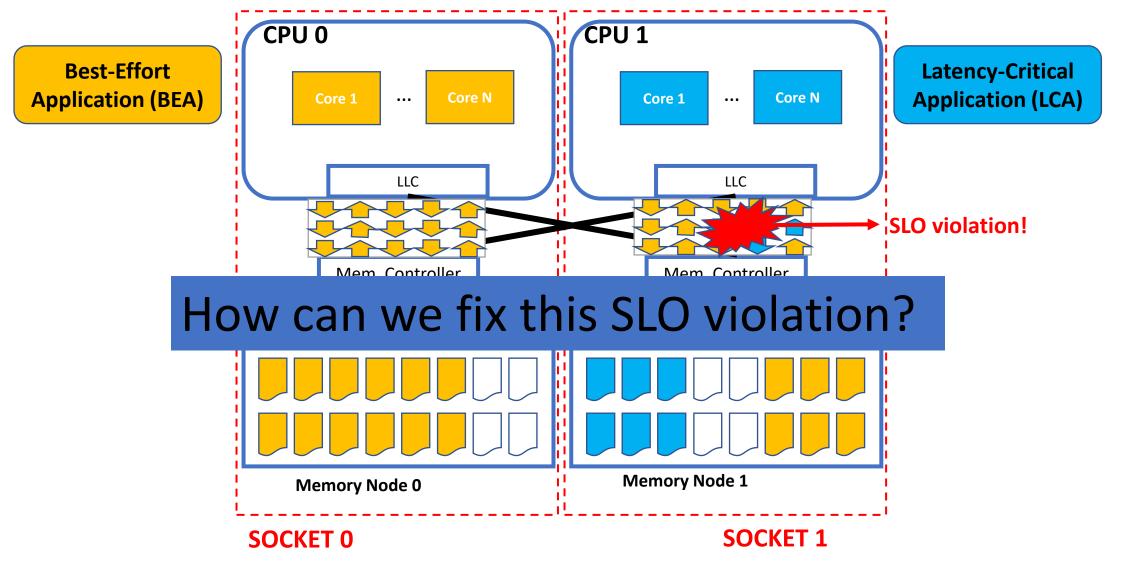


## Existing Solutions (2/2)

Disallows cross-socket sharing of memory

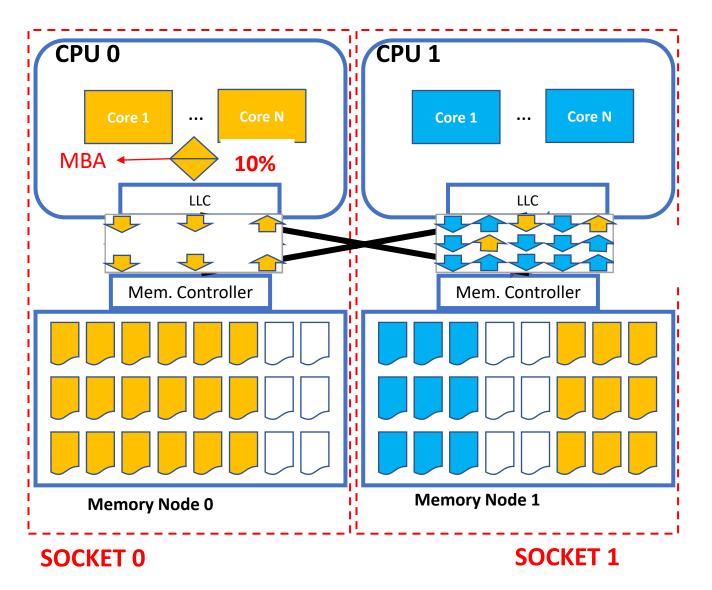
CPU 1 CPU 0 **Best-Effort Latency-Critical Application (BEA)** Core 1 **Core N Application (LCA)** LLC LLC Mem. Controller Mem. Controller **Memory Node 1 Memory Node 0 SOCKET 1 SOCKET 0** 

#### Cross-socket QoS-Aware Memory Bandwidth Allocation



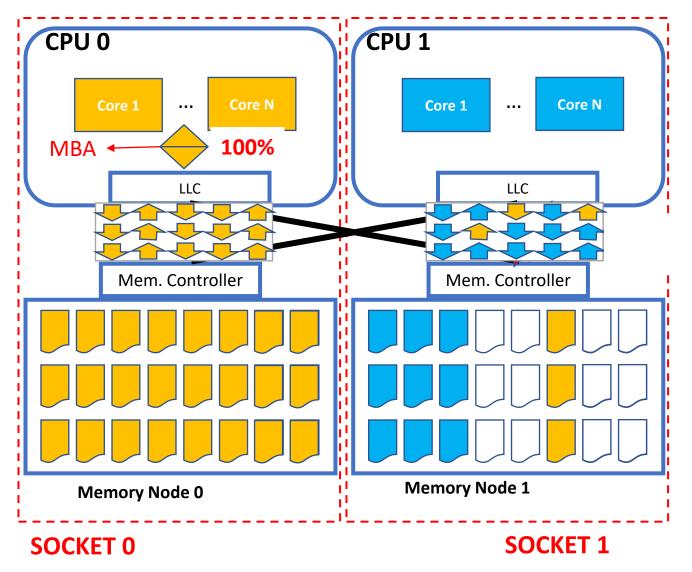
#### Fixing SLO Violation (1/2)

- Possibility 1: Intel Memory Bandwidth Allocation (MBA)
  - Can fix SLO violations almost instantaneously
  - Has a relevant cost on the performance of BEAs



#### Fixing SLO Violation (2/2)

- Possibility 2: Page Migration (pgm)
  - Can adjust memory bandwidth on a permemory node granularity
  - Page migration is slow but more efficient for BEA



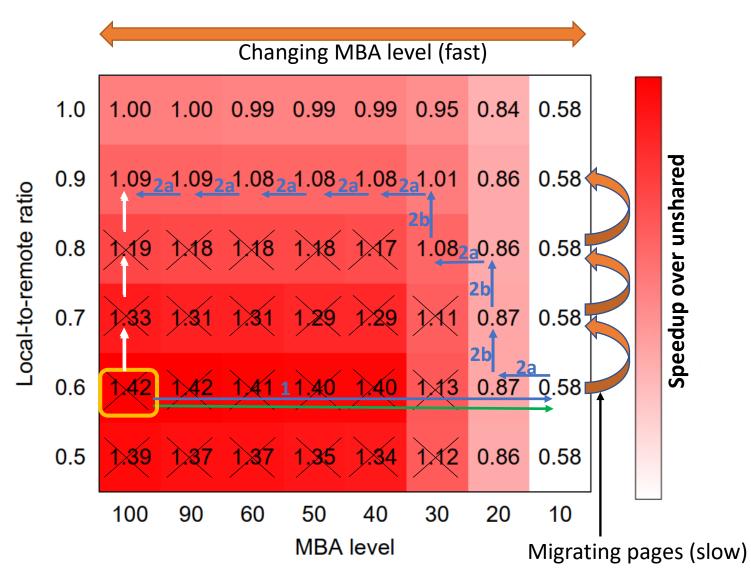
#### Contributions

- Study the hardware mechanism of Intel for memory bandwidth allocation (MBA) in a multi-socket scenario.
  - MBA unnecessarily reduces the throughput of BEAs by considerable margins

- We propose BALM, a novel QoS-aware memory bandwidth allocation technique for cross-socket sharing of memory in multi-socket architectures.
  - MBA with a novel cross-socket page migration scheme to obtain the best of both mechanisms

## Mitigating SLO Violations with BALM

- BALM uses MBA and page migration together as a 2-dimensional allocation mechanism
- BALM fixes SLO violations in 2 steps:
  - Set MBA to its most restrictive level
  - Apply incremental page migration
    - Gradually release MBA throttling
- BALM is expected to be:
  - i. As quick as MBA in fixing SLO violations
  - ii. Converge to valid optimal configuration as page migration



## Evaluation: Questions (1/4)

- 1. What performance advantage does BALM bring to memory-intensive BEAs on dual-socket NUMA systems?
- 2. How effective is BALM in fixing SLO violations?

- Compared BALM with:
  - MBA
  - Page migration (pgm)
  - Unshared

## Evaluation: Methodology (2/4)

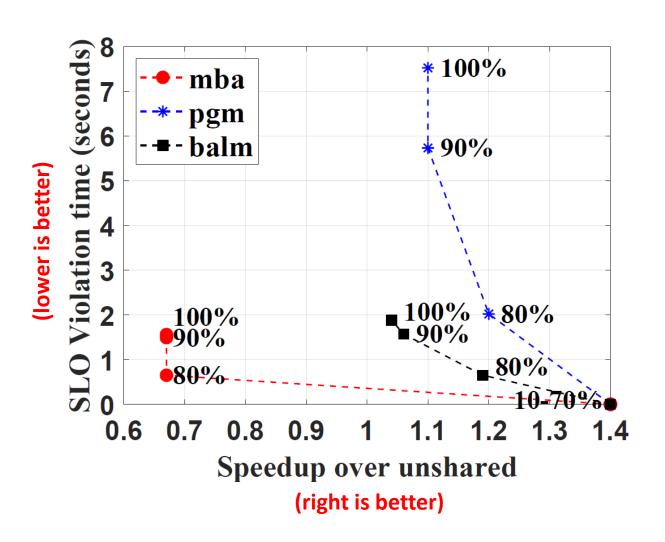
- BEAs: multithreaded benchmarks from PARSEC, SPLASH, NAS
- LCAs: Memcached

- Machines: dual-socket system
  - Intel Xeon Silver 4114 CPU, 2 NUMA nodes, 10 cores per node, supports MBA

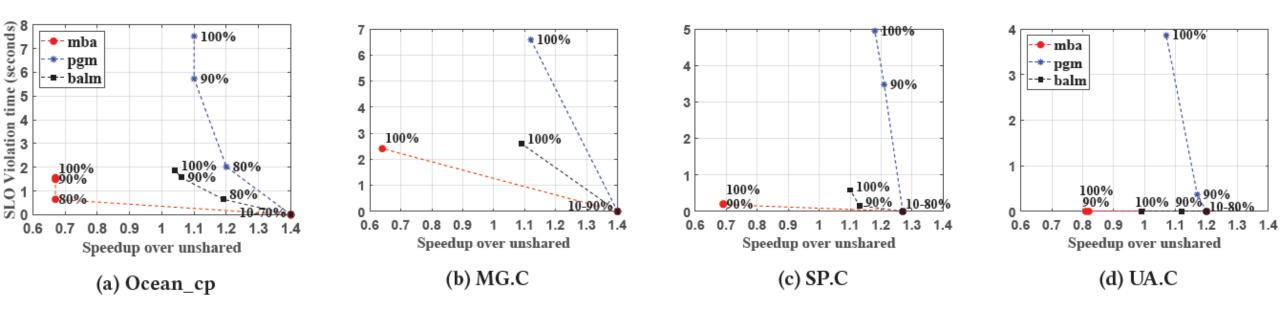
#### Execution Scenario

Socket 0	Socket 1
Memcached (4 threads)	OC/MG/SP/UA (8 threads)

## Evaluation: Results (Memcached vs. OC) (3/4)



## Evaluation: Results (Memcached vs. all) (4/4)



#### Conclusion

 SoTA QoS-aware resource allocation systems need to be generalized to allow cross-socket sharing of memory bandwidth

 BALM can safeguard the LCAs with marginal SLO violation windows, while delivering up to 87% throughput gains to bandwidth-intensive BEAs